

NaTT in CoCo 2022

Akihisa Yamada @ AIST

NaTT in CoCo

- NaTT a **termination prover** for **plain** term rewriting
- In CoCo 2021:
 - participated **INF** category with an easy checker [[Sternagel & Yamada '19](#)]
 - ignores conditions
 - weak but fast
- In CoCo 2022:
 - Generalized reduction pairs for INF [[IJCAR '22](#)]
 - don't ignore conditions
 - not too weak, but slow

Orderings for non-reachability

reduction pair [ver. Y, CADE '21]:

- \succcurlyeq and \succ are closed under substitutions
- \succcurlyeq is closed under contexts
- \succcurlyeq is a quasi-order
- \succ is well-founded (order)
- $\succcurlyeq \cdot \succ \sqsubseteq \succ$ **and** $\succ \cdot \succcurlyeq \sqsubseteq \succ$
- $\succ \sqsubseteq \succcurlyeq$

Orderings for non-reachability

rewrite pair [Y, IJCAR '22]:

- \succcurlyeq and \succ are closed under substitutions
- \succcurlyeq is closed under contexts
- \succcurlyeq is a quasi-order
- \succ is **irreflexive** (order)
- $\succcurlyeq \cdot \succ \sqsubseteq \succ$ and $\succ \cdot \succcurlyeq \sqsubseteq \succ$
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Theorem:

$s \rightarrow t$ is \mathcal{R} -unsat **if** there's a rewrite pair $\langle \succcurlyeq, \succ \rangle$ s.t. $\mathcal{R} \subseteq \succcurlyeq$ and $s < t$

Orderings for non-reachability

co-rewrite pair [Y, IJCAR '22]:

- \succcurlyeq and $<$ are closed under substitutions
- \succcurlyeq is closed under contexts
- \succcurlyeq is a quasi-order
- ($<$ is irreflexive)
- $\succcurlyeq \cap < = \emptyset$

Theorem:

$s \rightarrow t$ is \mathcal{R} -unsat **iff** there's a co-rewrite pair $\langle \succcurlyeq, < \rangle$ s.t. $\mathcal{R} \subseteq \succcurlyeq$ and $s < t$

Corollary:

$s \rightarrow t$ is \mathcal{R} -unsat **iff** \mathcal{R} has a quasi-model $\langle \mathcal{A}, \succeq \rangle$ s.t. $\mathcal{A} \models s \not\succeq t$

Co-WPO

$s \sqsupseteq_{\text{WPO}} t$

1. $\mathcal{A} \models s > t$ or
2. $\mathcal{A} \models s \geq t$ and
 - a. $\exists i \in \pi(f). s_i \sqsupseteq_{\text{WPO}} t$; or
 - b. $\forall j \in \pi(g). s \sqsupseteq_{\text{WPO}} t_j$ and
 - i. $f > g$ or
 - ii. $f \geq g$ and $\pi_f(s_1, \dots, s_n) \sqsupseteq_{\text{WPO}}^{\text{lex}} \pi_g(t_1, \dots, t_m)$;
 - c. $s = t \in \mathcal{V}$

$t \sqsupseteq_{\overline{\text{WPO}}} s$

1. $\mathcal{A} \models t \not\leq s$ or
2. $\mathcal{A} \models t \not\leq s$ and
 - a. $\exists j \in \pi(g). t_j \sqsupseteq_{\overline{\text{WPO}}} s$; or
 - b. $\forall i \in \pi(f). t \sqsupseteq_{\overline{\text{WPO}}} s_i$ and
 - i. $g \not\leq f$ or
 - ii. $g \not\leq f$ and $\pi_g(t_1, \dots, t_m) \sqsupseteq_{\overline{\text{WPO}}}^{\text{lex}} \pi_f(s_1, \dots, s_n)$;

Theorem [Y, IJCAR '22]:

$\langle \sqsupseteq_{\text{WPO}}, \sqsupseteq_{\overline{\text{WPO}}} \rangle$ is a co-rewrite pair (under mild condition).