

# infChecker at the 2020 Confluence Competition\*

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## 1 Overview

infChecker 1.0 is a tool for checking (*in*)feasibility of goals  $\mathcal{G} = \{F_i\}_{i=1}^m$  where  $F_i = (s_{ij} \bowtie_{ij} t_{ij})_{i=1}^{n_i}$  and  $\bowtie_{ij} \in \{\rightarrow, \rightarrow^*, \rightarrow^+, \leftrightarrow, \leftrightarrow^*, \leftrightarrow^+, \triangleright, \triangleright^+, \downarrow, \downarrow^+, \leftrightarrow, \leftrightarrow^*, \leftrightarrow^+, \leftrightarrow^*\}$  where predicates  $\bowtie_{ij}$  represent binary relations on terms (most of them well-known or easy generalizations of well-known relations) defined by provability of goals  $s \bowtie_{ij} t$  with respect to a *first-order theories*  $\text{Th}_{\bowtie_{ij}}$  [2, 4]. The tool is available here: <http://zenon.dsic.upv.es/infChecker/>. It is written in Haskell and provides a first implementation of the *Feasibility Framework* [2], where three *processors* have been implemented:

- $\text{P}^{\text{Sat}}$  integrates the satisfiability approach described in [3] to prove infeasibility. In infChecker, we use the model generators AGES [1] and Mace4 [6] to find a proof.
- $\text{P}^{\text{Prov}}$  integrates the logic-based approach to program analysis described in [3] to prove feasibility by theorem proving. In infChecker, we use the theorem prover Prover9 [6].
- $\text{P}^{\text{NC}}$  adapt the processor that narrow conditions in the 2D DP framework for proving operational termination of CTRs [5] to be used with feasibility sequences.

Our proof strategy is: (1) first, we try to prove feasibility using  $\text{P}^{\text{Prov}}$ ; (2) if  $\text{P}^{\text{Prov}}$  fails, we apply  $\text{P}^{\text{Sat}}$ ; (3) if  $\text{P}^{\text{Sat}}$  fails, we apply  $\text{P}^{\text{NC}}$ ; (4) if  $\text{P}^{\text{NC}}$  succeeds and modifies the feasibility sequence, we go to (2), otherwise we return MAYBE.

## References

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