

AGCP

System Description for CoCo 2016

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AGCP (Automated Ground Confluence Prover)

A ground confluence prover for many-sorted TRSs

- Available from

<http://www.nue.ie.niigata-u.ac.jp/tools/agcp/>

- SML/NJ is used for the implementation. Some program code are incorporated from confluence prover ACP [Aoto et al., 2009], and an inductive theorem prover [Aoto, 2008].

- entrant of **Demo(GCR)** category

- **Criteria:**

(1) \mathcal{R} is GCR if $\mathcal{S} \subseteq \mathcal{R}$ is GCR and $\mathcal{R} \setminus \mathcal{S}$ is a set of inductive theorems of \mathcal{S} .

(2) terminating \mathcal{S} is GCR if all critical pairs of \mathcal{S} are bounded ground convertible.

Terms s and t are *bounded ground convertible* wrt \approx if $s\sigma_g \overset{*}{\leftrightarrow}_{\approx \mathcal{R}} t\sigma_g$ holds for any ground substitution σ_g , where $s_g \overset{*}{\leftrightarrow}_{\approx \mathcal{R}} t_g$ if $s_g = u_0 \leftrightarrow_{\mathcal{R}} u_1 \leftrightarrow_{\mathcal{R}} \dots \leftrightarrow_{\mathcal{R}} u_n = t_g$ such that for any i , either $s \approx u_i$ or $t \approx u_i$.

- **Method:**

an extension of **rewriting induction** is used for proving both (1) and (2).

Rewriting Induction for GCR

Suppose \mathcal{R} is terminating and quasi-reducible. Let \succsim be a reduction quasi-order with $\mathcal{R} \subseteq \succsim$. Then, \mathcal{R} : GCR if $\langle \text{CP}(\mathcal{R}), \emptyset \rangle \overset{*}{\rightsquigarrow} \langle \emptyset, H \rangle$.

$$\frac{\langle E \uplus \{s \doteq t\}, H \rangle}{\langle E \cup \{s'_i \doteq t_i\}_i, H \cup \{s \rightarrow t\} \rangle} \quad \begin{array}{l} u \in \mathcal{B}(s), \\ \{s_i \rightarrow t_i\}_i = \text{Expd}_u(s, t), \\ s_i \overset{*}{\rightarrow}_{H \cup \text{inv}(H)} \tilde{s}'_i \end{array}$$

$$\frac{\langle E \uplus \{s \doteq t\}, H \rangle}{\langle E \cup \{s' \doteq t\}, H \rangle} s \rightarrow_{\mathcal{R} \cup H} \circ \overset{*}{\rightarrow}_{H \cup \text{inv}(H)} \tilde{s}'$$

$$\frac{\langle E \uplus \{s \doteq t\}, H \rangle}{\langle E, H \rangle} s \overset{=}{\leftrightarrow}_H t$$